

CHAPTER THREE

Section 3.1E

1.a. $2^1 = 2$

c. $2^2 = 4$

2.a. \downarrow

E	$\sim \sim (E \ \& \ \sim E)$
T	F T T F F T
F	F T F F T F

c. \downarrow

A	J	$A \equiv [J \equiv (A \equiv J)]$
T	T	T T T T T T T
T	F	T T F T T F F
F	T	F T T F F F T
F	F	F T F F F T F

e. \downarrow

A	H	J	$[\sim A \vee (H \supset J)] \supset (A \vee J)$
T	T	T	F T T T T T T T T T
T	T	F	F T F T F F T T T F
T	F	T	F T T F T T T T T T
T	F	F	F T T F T F T T T F
F	T	T	T F T T T T T F T T
F	T	F	T F T T F F F F F F
F	F	T	T F T F T T T F T T
F	F	F	T F T F T F F F F F

g. \downarrow

A	B	$\sim (A \vee B) \supset (\sim A \vee \sim B)$
T	T	F T T T T F T F F T
T	F	F T T F T F T T F
F	T	F F T T T T F T F T
F	F	T F F F T T F T T F

i. \downarrow

B	E	H	$\sim (E \ \& \ [H \supset (B \ \& \ E)])$
T	T	T	F T T T T T T T
T	T	F	F T T F T T T T
T	F	T	T F F T F T F F
T	F	F	T F F F T T F F
F	T	T	T T F T F F F T
F	T	F	F T T F T F F T
F	F	T	T F F T F F F F
F	F	F	T F F F T F F F

k.

D	E	F	\downarrow												
$\sim [D \ \& \ (E \ \vee \ F)]$			$\equiv \ [\sim D \ \& \ (E \ \& \ F)]$												
T	T	T	F	T	T	T	T	T	T	F	T	F	T	T	T
T	T	F	F	T	T	T	T	F	T	F	T	F	F	T	F
T	F	T	F	T	T	F	T	T	T	F	T	F	F	F	T
T	F	F	T	T	F	F	F	F	F	F	T	F	F	F	F
F	T	T	T	F	F	T	T	T	T	T	T	F	T	T	T
F	T	F	T	F	F	T	T	F	F	F	T	F	F	T	F
F	F	T	T	F	F	F	T	T	F	F	T	F	F	F	T
F	F	F	T	F	F	F	F	F	F	F	T	F	F	F	F

m.

A	H	J	\downarrow												
$(A \ \vee \ (\sim A \ \& \ (H \ \supset \ J)))$			$\supset \ (J \ \supset \ H)$												
T	T	T	T	T	F	T	F	T	T	T	T	T	T	T	T
T	T	F	T	T	F	T	F	T	F	F	T	F	T	T	T
T	F	T	T	T	F	T	F	F	T	T	F	T	F	F	F
T	F	F	T	T	F	T	F	F	T	F	T	F	T	F	F
F	T	T	F	T	T	F	T	T	T	T	T	T	T	T	T
F	T	F	F	F	T	F	F	T	F	F	T	F	T	T	T
F	F	T	F	T	T	F	T	T	F	T	T	F	T	F	F
F	F	F	F	T	T	F	T	T	F	T	F	T	F	T	F

3.a.

A	B	C	\downarrow								
$\sim [\sim A \ \vee \ (\sim C \ \vee \ \sim B)]$											
F	T	T	F	T	F	T	F	T	F	F	T

c.

A	B	C	\downarrow						
$(A \ \supset \ B) \ \vee \ (B \ \supset \ C)$									
F	T	T	F	T	T	T	T	T	T

e.

A	B	C	\downarrow						
$(A \ \equiv \ B) \ \vee \ (B \ \equiv \ C)$									
F	T	T	F	F	T	T	T	T	T

g.

A	B	C	\downarrow									
$\sim [B \ \supset \ (A \ \vee \ C)] \ \& \ \sim \sim B$												
F	T	T	F	T	T	F	T	T	F	T	F	T

i.

A	B	C	\downarrow												
$\sim [\sim (A \ \equiv \ \sim B) \ \equiv \ \sim A] \ \equiv \ (B \ \vee \ C)$															
F	T	T	T	F	F	T	F	T	F	T	F	T	T	T	T

4.a.

D	F	G	F	\vee	(G	\vee	D)
T	T	T	T	T	T	T	T
T	T	F	T	T	F	T	T
T	F	T	F	T	T	T	T
T	F	F	F	T	F	T	T
F	T	T	T	T	T	T	F
F	T	F	T	T	F	F	F
F	F	T	F	T	T	T	F
F	F	F	F	F	F	F	F

c.

D	F	G	[F	\vee	(G	\vee	D)]	&	(~	(F	\vee	G)	\vee	[~	(F	\vee	D)	\vee	~	(G	\vee	D)]				
T	T	T	T	T	T	T	F	F	T	T	T	F	F	T	T	T	F	F	T	T	T	F	F	T	T	T
T	T	F	T	T	F	T	F	F	T	T	F	F	F	T	T	T	F	F	T	T	T	F	F	F	T	T
T	F	T	F	T	T	T	F	F	F	T	T	F	F	F	T	T	F	F	F	T	T	F	F	T	T	T
T	F	F	F	T	F	T	T	T	F	F	F	F	T	F	F	T	T	F	F	F	T	T	F	F	F	T
F	T	T	T	T	T	F	F	F	T	T	T	F	F	T	T	F	F	T	T	F	F	F	T	T	F	T
F	T	F	T	T	F	F	F	T	F	T	T	F	T	F	T	T	F	T	T	F	T	T	F	F	F	F
F	F	T	F	T	T	F	T	F	F	T	T	T	T	F	F	F	T	T	F	F	F	T	F	T	T	F
F	F	F	F	F	F	F	F	T	F	F	F	T	T	F	F	F	T	T	F	F	F	T	T	F	F	F

e.

D	F	G	(F	&	G)	\vee	[(F	&	D)	\vee	(G	&	D)]
T	T	T	T	T	T	T	T	T	T	T	T	T	T
T	T	F	T	F	F	T	T	T	T	T	F	F	T
T	F	T	F	F	T	T	F	F	T	T	T	T	T
T	F	F	F	F	F	F	F	F	T	F	F	F	T
F	T	T	T	T	T	T	T	F	F	F	T	F	F
F	T	F	T	F	F	F	T	F	F	F	F	F	F
F	F	T	F	F	T	F	F	F	F	F	T	F	F
F	F	F	F	F	F	F	F	F	F	F	F	F	F

g.

D	F	G	[(F	&	G)	&	~	D]	\vee	[(F	&	D)	&	~	G]	\vee	[(G	&	D)	&	~	F]		
T	T	T	T	T	T	F	F	T	F	T	T	T	F	F	T	F	T	T	T	F	F	T		
T	T	F	T	F	F	F	F	T	T	T	T	T	F	T	T	F	F	T	F	F	T	F	F	T
T	F	T	F	F	T	F	F	T	F	F	T	F	F	T	T	T	T	T	F	T	F	F	T	F
T	F	F	F	F	F	F	F	T	F	F	T	F	F	T	F	F	T	F	T	F	F	T	F	T
F	T	T	T	T	T	T	T	F	T	F	F	F	F	T	F	F	T	F	F	F	F	T	F	T
F	T	F	T	F	F	F	T	F	T	F	F	F	T	F	F	F	F	F	F	F	F	T	F	T
F	F	T	F	F	T	F	T	F	F	F	F	F	F	T	F	F	T	F	F	F	T	F	T	F
F	F	F	F	F	F	F	T	F	F	F	F	T	F	F	F	F	F	F	F	T	F	T	F	T

5.a.

D	F	G		[F \vee (G \vee D)]	\supset	[F & (G & D)]
T	T	T		T	T	T
T	T	F		T	F	F
T	F	T		F	F	T
T	F	F		F	F	F
F	T	T		T	T	F
F	T	F		T	F	F
F	F	T		F	F	T
F	F	F		F	F	F

c.

D	F	G	S		S \supset [G \supset \sim (F \vee D)]
T	T	T	T		T
T	T	T	F		F
T	T	F	T		T
T	T	F	F		F
T	F	T	T		T
T	F	T	F		F
T	F	F	T		T
T	F	F	F		F
F	T	T	T		T
F	T	T	F		F
F	T	F	T		T
F	T	F	F		F
F	F	T	T		T
F	F	T	F		F
F	F	F	T		T
F	F	F	F		F

e.

D	P	S		D \equiv (P & S)
T	T	T		T
T	T	F		F
T	F	T		F
T	F	F		F
F	T	T		T
F	T	F		F
F	F	T		F
F	F	F		F

g.

D	F	G	P	R	P	\supset	(F	\supset	$[\sim$	(D	\vee	G)	&	R])
T	T	T	T	T	T	F	T	F	F	T	T	T	F	T
T	T	T	T	F	T	F	T	F	F	T	T	T	F	F
T	T	T	F	T	F	T	T	F	F	T	T	T	F	T
T	T	T	F	F	F	T	T	F	F	T	T	T	F	F
T	T	F	T	T	T	F	T	F	F	T	T	F	F	T
T	T	F	T	F	T	F	T	F	F	T	T	F	F	F
T	T	F	F	T	F	T	T	F	F	T	T	F	F	T
T	T	F	F	F	F	T	T	F	F	T	T	F	F	F
T	F	T	T	T	T	T	F	T	F	T	T	T	F	T
T	F	T	T	F	T	T	F	T	F	T	T	T	F	F
T	F	T	F	T	F	T	F	T	F	T	T	T	F	T
T	F	T	F	F	F	T	F	T	F	T	T	T	F	F
T	F	F	T	T	F	T	F	T	F	T	T	F	F	T
T	F	F	T	F	F	T	F	T	F	T	T	F	F	F
T	F	F	F	T	F	T	F	T	F	T	T	F	F	T
T	F	F	F	F	F	T	F	T	F	T	T	F	F	F
F	T	T	T	T	T	F	T	F	F	F	T	T	F	T
F	T	T	T	F	T	F	T	F	F	F	T	T	F	F
F	T	T	F	T	F	T	T	F	F	F	T	T	F	T
F	T	T	F	F	F	T	T	F	F	F	T	T	F	F
F	T	F	T	T	T	T	T	T	T	F	F	F	T	T
F	T	F	T	F	T	F	T	F	T	F	F	F	F	F
F	T	F	F	T	F	T	T	T	T	F	F	F	T	T
F	T	F	F	F	F	T	T	F	T	F	F	F	F	F
F	T	F	F	F	F	T	T	F	T	F	F	F	F	F
F	F	T	T	T	T	T	F	T	F	F	T	T	F	T
F	F	T	T	F	T	T	F	T	F	F	T	T	F	F
F	F	T	F	T	F	T	F	T	F	F	T	T	F	T
F	F	T	F	F	F	T	F	T	F	F	T	T	F	F
F	F	F	T	T	T	T	F	T	T	F	F	F	T	T
F	F	F	T	F	T	T	F	T	T	F	F	F	F	F
F	F	F	F	T	F	T	F	T	T	F	F	F	T	T
F	F	F	F	F	F	T	F	T	T	F	F	F	F	F

Section 3.2E

1.a. Truth-functionally indeterminate

A	\sim A	\supset	A
T	F	T	T
F	T	F	F

c. Truth-functionally true

$$\downarrow$$

A	$(A \equiv \sim A) \supset \sim(A \equiv \sim A)$								
T	T	F	F	T	T	T	F	F	T
F	F	F	T	F	T	F	F	T	F

e. Truth-functionally indeterminate

$$\downarrow$$

B	D	$(\sim B \ \& \ \sim D) \vee \sim(B \vee D)$							
T	T	F	T	F	F	T	T	T	T
T	F	F	T	F	F	T	T	F	F
F	T	T	F	F	F	F	F	T	T
F	F	T	F	T	F	T	T	F	F

g. Truth-functionally indeterminate

$$\downarrow$$

A	B	C	$[(A \vee B) \ \& \ (A \vee C)] \supset \sim(B \ \& \ C)$								
T	T	T	T	T	T	T	T	F	F	T	T
T	T	F	T	T	T	T	F	T	T	F	F
T	F	T	T	T	F	T	T	T	T	F	F
T	F	F	T	T	F	T	F	T	T	F	F
F	T	T	F	T	T	F	T	F	F	T	T
F	T	F	F	T	F	F	F	T	T	F	F
F	F	T	F	F	F	F	T	T	T	F	T
F	F	F	F	F	F	F	F	T	T	F	F

i. Truth-functionally true

$$\downarrow$$

J	K	$(J \vee \sim K) \equiv \sim\sim(K \supset J)$							
T	T	T	T	F	T	T	F	T	T
T	F	T	T	T	F	T	F	F	T
F	T	F	F	F	T	F	T	F	F
F	F	F	T	T	F	T	F	F	T

k. Truth-functionally true

$$\downarrow$$

A	D	$[(A \vee \sim D) \ \& \ \sim(A \ \& \ D)] \supset \sim D$							
T	T	T	T	F	F	F	T	T	T
T	F	T	T	T	F	T	T	F	F
F	T	F	F	F	F	T	F	F	T
F	F	F	T	T	F	T	F	F	T

2.a. Not truth-functionally true

$$\begin{array}{c|c} & \downarrow \\ \hline F & H \\ \hline F & F \end{array} \quad \begin{array}{c} (F \vee H) \vee (\sim F \equiv H) \\ \hline F \quad F \quad F \quad F \quad T \quad F \quad F \quad F \end{array}$$

c. Truth-functionally true

$$\begin{array}{c|c} & \downarrow \\ \hline A & B & C \\ \hline T & T & T \\ T & T & F \\ T & F & T \\ T & F & F \\ F & T & T \\ F & T & F \\ F & F & T \\ F & F & F \end{array} \quad \begin{array}{c} \sim A \supset [(B \& A) \supset C] \\ \hline F \quad T \quad T \quad T \quad T \quad T \quad T \quad T \\ F \quad T \quad T \quad T \quad T \quad T \quad F \quad F \\ F \quad T \quad T \quad F \quad F \quad T \quad T \quad T \\ F \quad T \quad T \quad F \quad F \quad T \quad T \quad F \\ T \quad F \quad T \quad T \quad F \quad F \quad T \quad T \\ T \quad F \quad T \quad T \quad F \quad F \quad T \quad F \\ T \quad F \quad T \quad T \quad F \quad F \quad T \quad F \\ T \quad F \quad T \quad F \quad F \quad F \quad T \quad F \\ T \quad F \quad T \quad F \quad F \quad F \quad T \quad T \\ T \quad F \quad T \quad F \quad F \quad F \quad T \quad F \end{array}$$

e. Truth-functionally true

$$\begin{array}{c|c} & \downarrow \\ \hline C \\ \hline T \\ F \end{array} \quad \begin{array}{c} [(C \vee \sim C) \supset C] \supset C \\ \hline T \quad T \quad F \quad T \quad T \quad T \quad T \\ F \quad T \quad T \quad F \quad F \quad F \quad T \quad F \end{array}$$

3.a. Truth-functionally false

$$\begin{array}{c|c} & \downarrow \\ \hline B & D \\ \hline T & T \\ T & F \\ F & T \\ F & F \end{array} \quad \begin{array}{c} (B \equiv D) \& (B \equiv \sim D) \\ \hline T \quad T \quad T \quad F \quad T \quad F \quad F \quad T \\ T \quad F \quad F \quad F \quad T \quad T \quad T \quad F \\ F \quad F \quad T \quad F \quad F \quad T \quad F \quad T \\ F \quad T \quad F \quad F \quad F \quad F \quad T \quad F \end{array}$$

c. Not truth-functionally false

$$\begin{array}{c|c} & \downarrow \\ \hline A & B \\ \hline T & T \\ T & T \end{array} \quad \begin{array}{c} A \equiv (B \equiv A) \\ \hline T \quad T \quad T \quad T \quad T \end{array}$$

e. Not truth-functionally false

$$\begin{array}{c|c} & \downarrow \\ \hline C & D \\ \hline F & T \end{array} \quad \begin{array}{c} [(C \vee D) \equiv C] \supset \sim C \\ \hline F \quad T \quad T \quad F \quad F \quad T \quad T \quad F \end{array}$$

4.a. False. For example, while ' $(A \supset A)$ ' is truth-functionally true, ' $(A \supset A) \& A$ ' is not.

c. True. There cannot be any truth-value assignment on which the antecedent is true and the consequent false because there is no truth-value assignment on which the consequent is false.

e. False. For example, although ' $(A \& \sim A)$ ' is truth-functionally false, ' $C \vee (A \& \sim A)$ ' is not.

g. True. Since a sentence $\sim \mathbf{P}$ is false on a truth-value assignment if and only if \mathbf{P} is true on the truth-value assignment, \mathbf{P} is truth-functionally true if and only if $\sim \mathbf{P}$ is truth-functionally false.

i. False. For example, ' $(A \vee \sim A)$ ' is truth-functionally true, but ' $(A \vee \sim A) \supset B$ ' is truth-functionally indeterminate.

5.a. On every truth-value assignment, \mathbf{P} is true and \mathbf{Q} is false. Hence $\mathbf{P} \equiv \mathbf{Q}$ is false on every truth-value assignment. Therefore $\mathbf{P} \equiv \mathbf{Q}$ is truth-functionally false.

c. No. Both ' A ' and ' $\sim A$ ' are truth-functionally indeterminate, but ' $A \vee \sim A$ ' is truth-functionally true.

Section 3.3E

1.a. Not truth-functionally equivalent

A	B	\downarrow $\sim (A \& B)$	\downarrow $\sim (A \vee B)$
T	T	F	F
T	F	T	T
F	T	T	T
F	F	T	T

c. Truth-functionally equivalent

H	K	\downarrow $K \equiv H$	\downarrow $\sim K \equiv \sim H$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	T	T

e. Truth-functionally equivalent

F	G	\downarrow $(G \supset F) \supset (F \supset G)$	\downarrow $(G \equiv F) \vee (\sim F \vee G)$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

g. Not truth-functionally equivalent

H	J	K	\downarrow $\sim (H \ \& \ J) \equiv (J \equiv \sim K)$						\downarrow $(H \ \& \ J) \supset \sim K$					
T	T	T	F	T	T	T	T	F	FT	T	T	T	F	FT
T	T	F	F	T	T	F	T	T	TF	T	T	T	T	TF
T	F	T	T	T	F	F	T	F	FT	T	F	F	T	FT
T	F	F	T	T	F	F	F	F	TF	T	F	F	T	TF
F	T	T	T	F	F	T	F	T	FT	F	F	T	T	FT
F	T	F	T	F	F	T	T	T	TF	F	F	T	T	TF
F	F	T	T	F	F	T	F	T	FT	F	F	F	T	FT
F	F	F	T	F	F	F	F	F	TF	F	F	F	T	TF

i. Not truth-functionally equivalent

A	C	D	\downarrow $[A \vee \sim (D \ \& \ C)] \supset \sim D$						\downarrow $[D \vee \sim (A \ \& \ C)] \supset \sim A$								
T	T	T	T	T	F	T	T	F	FT	T	T	F	T	T	T	F	FT
T	T	F	T	T	T	F	F	T	T	TF	F	F	F	T	T	T	T
T	F	T	T	T	T	F	F	F	FT	T	T	T	T	F	F	F	FT
T	F	F	T	T	T	F	F	F	T	TF	F	T	T	T	F	F	F
F	T	T	F	F	F	T	T	T	T	FT	T	T	T	F	F	T	TF
F	T	F	F	T	T	F	F	T	TF	F	T	T	F	F	T	T	TF
F	F	T	F	T	T	F	F	F	FT	T	T	T	F	F	F	T	TF
F	F	F	F	T	T	F	F	F	T	TF	F	T	T	F	F	F	T

k. Not truth-functionally equivalent

F	G	H	\downarrow $F \vee \sim (G \vee \sim H)$						\downarrow $(H \equiv \sim F) \vee G$				
T	T	T	T	T	F	T	T	FT	T	F	FT	T	T
T	T	F	T	T	F	T	T	TF	F	T	FT	T	T
T	F	T	T	T	T	F	F	FT	T	F	FT	F	F
T	F	F	T	T	F	F	T	TF	F	T	FT	T	F
F	T	T	F	F	F	T	T	FT	T	T	TF	T	T
F	T	F	F	F	F	T	T	TF	F	F	TF	T	T
F	F	T	F	T	T	F	F	FT	T	T	TF	T	F
F	F	F	F	F	F	F	T	TF	F	F	TF	F	F

2.a. Truth-functionally equivalent

G	H	\downarrow $G \vee H$				\downarrow $\sim G \supset H$			
T	T	T	T	T	FT	T	T		
T	F	T	T	F	FT	T	F		
F	T	F	T	T	TF	T	T		
F	F	F	F	F	TF	F	F		

c. Truth-functionally equivalent

A	D	$(D \equiv A)$	$\&$	D	$\&$	A
T	T	T	T	T	T	T
T	F	F	F	T	F	T
F	T	T	F	F	T	F
F	F	F	T	F	F	F

e. Not truth-functionally equivalent

A	A	\equiv	$(\sim A \equiv A)$	$\sim (A \supset \sim A)$
T	T	F	F	T
T	F	F	T	F
F	T	T	F	T
F	F	T	T	F

3.a. Not truth-functionally equivalent

C: The sky clouds over.

N: The night will be clear.

M: The moon will shine brightly.

C	M	N	C	\vee	(N	$\&$	M)	M	\equiv	(N	$\&$	\sim	C)
T	T	T	T	T	T	T	T	T	F	T	F	F	T
T	T	F	T	T	F	F	T	T	F	F	F	F	T
T	F	T	T	T	T	F	F	F	T	F	F	F	T
T	F	F	T	T	F	F	F	F	T	F	F	F	T
F	T	T	F	T	T	T	T	T	T	T	T	T	F
F	T	F	F	F	F	F	T	T	F	F	F	T	F
F	F	T	F	F	T	F	F	F	F	T	T	T	F
F	F	F	F	F	F	F	F	F	T	F	F	T	F

c. Truth-functionally equivalent

D: The *Daily Herald* reports on our antics.

A: Our antics are effective.

A	D	D	\supset	A	$\sim A$	\supset	$\sim D$
T	T	T	T	T	F	T	F
T	F	F	T	T	F	T	T
F	T	T	F	F	T	F	F
F	F	F	T	F	T	T	F

e. Not truth-functionally equivalent

M: Mary met Tom.

L: Mary liked Tom.

G: Mary asked George to the movies.

G	L	M	(M & L)	\downarrow $\supset \sim G$	(M & $\sim L$)	\downarrow $\supset G$
T	T	T	T	F	T	T
T	T	F	F	T	F	T
T	F	T	F	T	T	T
T	F	F	F	T	F	T
F	T	T	T	T	F	F
F	T	F	F	T	F	F
F	F	T	T	T	T	F
F	F	F	F	T	F	F

4.a. Yes. **P** and **Q** have the same truth-value on every truth-value assignment. On every truth-value assignment on which they are both true, $\sim \mathbf{P}$ and $\sim \mathbf{Q}$ are both false, and on every truth-value assignment on which they are both false, $\sim \mathbf{P}$ and $\sim \mathbf{Q}$ are both true. It follows that $\sim \mathbf{P}$ and $\sim \mathbf{Q}$ are truth-functionally equivalent.

c. If **P** and **Q** are truth-functionally equivalent then they have the same truth-value on every truth-value assignment. On those assignments on which they are both true, the second disjunct of $\sim \mathbf{P} \vee \mathbf{Q}$ is true and so is the disjunction. On those assignments on which they are both false, the first disjunct of $\sim \mathbf{P} \vee \mathbf{Q}$ is true and so is the disjunction. So $\sim \mathbf{P} \vee \mathbf{Q}$ is true on every truth-value assignment.

Section 3.4E

1.a. Truth-functionally consistent

A	B	C	\downarrow A \supset B	\downarrow B \supset C	\downarrow A \supset C
T	T	T	T	T	T
T	T	F	T	F	F
T	F	T	F	T	T
T	F	F	F	F	F
F	T	T	F	T	T
F	T	F	F	F	F
F	F	T	F	T	T
F	F	F	F	F	F

c. Truth-functionally inconsistent

H	J	L	\downarrow $\sim [J \vee (H \supset L)]$	\downarrow $L \equiv (\sim J \vee \sim H)$	\downarrow $H \equiv (J \vee L)$
T	T	T	F	T	T
T	T	F	F	T	T
T	F	T	F	T	T
T	F	F	T	F	F
F	T	T	F	T	T
F	T	F	F	T	T
F	F	T	F	T	T
F	F	F	F	T	T

e. Truth-functionally inconsistent

H	J	\downarrow $(J \supset J) \supset H$	\downarrow $\sim J$	\downarrow $\sim H$
T	T	T	F	F
T	F	F	T	F
F	T	T	F	T
F	F	F	T	T

g. Truth-functionally consistent

A	B	C	\downarrow A	\downarrow B	\downarrow C
T	T	T	T	T	T
T	T	F	T	T	F
T	F	T	T	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	T	F	F	T	F
F	F	T	F	F	T
F	F	F	F	F	F

i. Truth-functionally consistent

A	B	C	\downarrow $(A \& B) \vee (C \supset B)$	\downarrow $\sim A$	\downarrow $\sim B$
T	T	T	T	F	F
T	T	F	T	F	F
T	F	T	F	F	T
T	F	F	T	F	T
F	T	T	F	T	F
F	T	F	F	T	F
F	F	T	F	T	T
F	F	F	F	T	T

2.a. Truth-functionally consistent

	\downarrow		\downarrow
B D E		$B \supset (D \supset E)$	$\sim D \ \& \ B$
T F T		T T F T T	T F T T

c. Truth-functionally consistent

	\downarrow		\downarrow
F J K		$F \supset (J \vee K)$	$F \equiv \sim J$
T F T		T T F T T	T T T F

e. Truth-functionally consistent

	\downarrow		\downarrow
A B		$(A \supset B) \equiv (\sim B \vee B)$	A
T T		T T T T F T T T	T

3.a. Truth-functionally inconsistent

S: Space is infinitely divisible.

Z: Zeno's paradoxes are compelling.

C: Zeno's paradoxes are convincing.

	\downarrow		\downarrow		\downarrow
C S Z		$S \supset Z$	$\sim (C \vee Z)$		S
T T T		T T T	F T T T		T
T T F		T F F	F T T F		T
T F T		F T T	F T T T		F
T F F		F T F	F T T F		F
F T T		T T T	F F T T		T
F T F		T F F	T F F F		T
F F T		F T T	F F T T		F
F F F		F T F	T F F F		F

c. Truth-functionally consistent

E: Eugene O'Neill was an alcoholic.

P: Eugene O'Neill's plays show that he was an alcoholic.

I: *The Iceman Cometh* must have been written by a teetotaler.

F: Eugene O'Neill was a fake.

				↓	↓	↓	↓			
E	F	I	P	E	P	I	E	∨	F	
T	T	T	T	T	T	T	T	T	T	T
T	T	T	F	T	F	T	T	T	T	T
T	T	F	T	T	T	F	T	T	T	T
T	T	F	F	T	F	F	T	T	T	T
T	F	T	T	T	T	T	T	T	F	F
T	F	T	F	T	F	T	T	T	F	F
T	F	F	T	T	T	F	T	T	F	F
T	F	F	F	T	F	F	T	T	F	F
F	T	T	T	F	T	T	F	T	T	T
F	T	T	F	F	F	T	F	T	T	T
F	T	F	T	F	T	F	F	F	T	T
F	T	F	F	F	F	F	F	F	T	T
F	F	T	T	F	T	T	F	F	F	F
F	F	T	F	F	F	T	F	F	F	F
F	F	F	T	F	T	F	F	F	F	F
F	F	F	F	F	F	F	F	F	F	F

e. Truth-functionally consistent

R: The Red Sox will win next Sunday.

J: Joan bet \$5.00.

E: Joan will buy Ed a hamburger.

			↓				↓					
E	J	R	R	⊃	(J	⊃	E)	~ R	&	~ E		
T	T	T	T	T	T	T	T	F	T	F	F	T
T	T	F	F	T	T	F	T	T	F	F	F	F
T	F	T	T	T	T	F	T	T	F	T	F	F
T	F	F	F	T	F	F	T	T	T	F	F	F
F	T	T	T	F	T	F	F	F	T	F	F	T
F	T	F	F	T	T	F	F	F	T	F	T	F
F	F	T	T	T	T	F	T	F	F	T	F	T
F	F	F	F	F	T	F	T	F	F	T	F	T

4.a. First assume that $\{\mathbf{P}\}$ is truth-functionally inconsistent. Then, since \mathbf{P} is the only member of $\{\mathbf{P}\}$, there is no truth-value assignment on which \mathbf{P} is true; so \mathbf{P} is false on every truth-value assignment. But then $\sim \mathbf{P}$ is true on every truth-value assignment, and so $\sim \mathbf{P}$ is truth-functionally true.

Now assume that $\sim \mathbf{P}$ is truth-functionally true. Then $\sim \mathbf{P}$ is true on every truth-value assignment, and so \mathbf{P} is false on every truth-value assignment. But then there is no truth-value assignment on which \mathbf{P} , the only member of $\{\mathbf{P}\}$, is true, and so the set is truth-functionally inconsistent.

c. No. For example, 'A' and ' $\sim A$ ' are both truth-functionally indeterminate, but $\{A, \sim A\}$ is truth-functionally inconsistent.

Section 3.5E

1.a. Truth-functionally valid

A	H	J	↓ A \supset (H & J)				↓ J \equiv H			↓ \sim J	↓ \sim A
T	T	T	T	T	T	T	T	T	T	FT	FT
T	T	F	T	F	T	F	F	F	T	TF	FT
T	F	T	T	F	F	F	T	T	F	FT	FT
T	F	F	T	F	F	F	F	F	F	TF	FT
F	T	T	F	T	T	T	T	T	T	FT	TF
F	T	F	F	T	T	F	F	F	T	TF	TF
F	F	T	F	T	F	F	T	T	F	FT	TF
F	F	F	F	T	F	F	F	F	T	TF	TF

c. Truth-functionally valid

A	D	G	↓ (D \equiv \sim G) & G				↓ (G \vee [(A \supset D) & A])				↓ \supset \sim D	↓ G \supset \sim D				
T	T	T	T	F	FT	F	T	T	T	T	T	F	FT	T	F	FT
T	T	F	T	T	TF	F	F	T	T	T	T	F	FT	F	T	FT
T	F	T	F	T	FT	T	T	T	F	F	F	T	TF	T	T	TF
T	F	F	F	F	TF	F	F	F	T	F	F	T	TF	F	T	TF
F	T	T	T	F	FT	F	T	T	F	F	F	F	FT	T	F	FT
F	T	F	T	T	TF	F	F	F	T	T	F	F	T	FT	F	FT
F	F	T	F	T	FT	T	T	T	F	T	F	F	T	TF	T	TF
F	F	F	F	F	TF	F	F	F	T	F	F	F	T	TF	F	TF

e. Truth-functionally valid

C	D	E	\downarrow (C \supset D) \supset (D \supset E)				\downarrow D	\downarrow C \supset E		
T	T	T	T	T	T	T	T	T	T	T
T	T	F	T	T	T	F	T	F	F	F
T	F	T	T	F	F	T	F	T	T	T
T	F	F	T	F	F	T	F	T	F	F
F	T	T	F	T	T	T	T	T	T	T
F	T	F	F	T	T	F	T	F	F	F
F	F	T	F	T	F	T	T	F	T	T
F	F	F	F	T	F	T	F	F	T	F

g. Truth-functionally valid

G	H	\downarrow (G \equiv H) \vee (\sim G \equiv H)				\downarrow (\sim G \equiv \sim H) \vee \sim (G \equiv H)			
T	T	T	T	T	F	T	F	T	T
T	F	T	F	F	T	F	T	T	F
F	T	F	F	T	T	F	T	T	T
F	F	F	T	F	T	F	T	F	F

i. Truth-functionally invalid

F	G	\downarrow $\sim\sim$ F \supset $\sim\sim$ G			\downarrow \sim G \supset \sim F			\downarrow G \supset F		
T	T	T	F	T	T	F	T	T	T	
T	F	T	F	F	T	F	F	T	T	
F	T	F	T	T	F	T	T	F	F	
F	F	F	T	F	T	F	F	F	F	

2.a. Truth-functionally valid

J	M	\downarrow (J \vee M) \supset \sim (J & M)				\downarrow M \equiv (M \supset J)				\downarrow M \supset J			
T	T	T	T	T	F	F	T	T	T	T	T	T	T
T	F	T	T	F	T	T	T	F	F	F	F	T	T
F	T	F	T	T	T	T	F	F	T	F	F	T	F
F	F	F	F	F	T	T	F	F	F	F	F	F	F

c. Truth-functionally valid

		↓					↓				↓		
A	B	$A \supset \sim A$			$(B \supset A) \supset B$				$A \equiv \sim B$				
T	T	T	F	F	T	T	T	T	T	T	F	F	F
T	F	T	F	F	F	T	T	F	F	T	T	T	F
F	T	F	T	T	T	F	F	T	T	F	T	F	T
F	F	F	T	T	F	T	F	F	F	F	F	F	T

e. Truth-functionally invalid

			↓							↓				↓		
A	B	C	$A \& \sim [(B \& C) \equiv (C \supset A)]$						$B \supset \sim B$			$\sim C \supset C$				
T	F	F	T	T	T	F	F	F	F	F	T	T	F	T	F	F

3.a. Truth-functionally valid

		↓
B	C	$(B \& C) \supset (B \vee C)$
T	T	T
T	F	T
F	T	T
F	F	T

c. Truth-functionally invalid

		↓
J	T	$[(J \supset T) \supset J] \& [(T \supset J) \supset T] \supset (\sim J \vee \sim T)$
T	T	F

e. Truth-functionally invalid

			↓
B	C	D	$[(B \& C) \& (B \vee D)] \supset D$
T	T	F	F

4.a. Truth-functionally invalid

S: 'Stern' means the same as 'star'.

N: 'Nacht' means the same as 'day'.

		↓	↓	↓
N	S	$N \supset S$	$\sim N$	$\sim S$
T	T	T	F	F
T	F	F	F	T
F	T	T	T	F
F	F	F	T	T

c. Truth-functionally valid

S: September has 30 days.

A: April has 30 days.

N: November has 30 days.

F: February has 40 days.

M: May has 30 days.

A	F	M	N	S	↓ S & (A & N)	(A ≡ ~ M)	↓ & (N ⊃ M)	↓ F
T	T	T	T	T	T	T	F	T
T	T	T	T	F	F	F	F	T
T	T	T	F	T	T	F	F	T
T	T	T	F	F	F	F	F	T
T	T	F	T	T	T	T	F	T
T	T	F	T	F	F	T	F	T
T	T	F	F	T	T	T	F	T
T	T	F	F	F	F	T	F	T
T	F	T	T	T	T	F	F	F
T	F	T	T	F	F	F	F	F
T	F	T	F	T	T	F	F	F
T	F	T	F	F	F	F	F	F
T	F	F	T	T	T	T	F	F
T	F	F	T	F	F	T	F	F
T	F	F	F	T	T	T	F	F
T	F	F	F	F	F	T	F	F
F	T	T	T	T	T	F	T	T
F	T	T	T	F	F	F	T	T
F	T	T	F	T	F	F	T	T
F	T	T	F	F	F	F	T	T
F	T	F	T	T	T	F	F	T
F	T	F	T	F	F	F	F	T
F	T	F	F	T	T	F	F	T
F	T	F	F	F	F	F	F	T
F	F	T	T	T	T	F	T	F
F	F	T	T	F	F	F	T	F
F	F	T	F	T	T	F	F	F
F	F	T	F	F	F	F	F	F
F	F	F	T	T	T	F	T	F
F	F	F	T	F	F	F	F	F
F	F	F	F	T	T	F	F	F
F	F	F	F	F	F	F	F	F

e. Truth-functionally valid

D: Computers can have desires.

E: Computers can have emotions.

T: Computers can think.

D	E	T	↓	T ≡ E	↓	E ⊃ D	↓	D ⊃ ~T	↓	~T
T	T	T	↓	T	↓	T	↓	T	↓	F
T	T	F	↓	F	↓	T	↓	T	↓	T
T	F	T	↓	T	↓	F	↓	F	↓	F
T	F	F	↓	F	↓	F	↓	T	↓	T
F	T	T	↓	T	↓	T	↓	F	↓	F
F	T	F	↓	F	↓	T	↓	F	↓	T
F	F	T	↓	T	↓	F	↓	T	↓	F
F	F	F	↓	F	↓	F	↓	F	↓	T

5.a. Suppose that the argument is truth-functionally valid. Then there is no truth-value assignment on which $\mathbf{P}_1, \dots, \mathbf{P}_n$ are all true and \mathbf{Q} is false. But, by the characteristic truth-table for ‘&’, the iterated conjunction $(\dots (\mathbf{P}_1 \& \mathbf{P}_2) \& \dots \mathbf{P}_n)$ has the truth-value \mathbf{T} on a truth-value assignment if and only if all of $\mathbf{P}_1, \dots, \mathbf{P}_n$ have the truth-value \mathbf{T} on that assignment. So, on our assumption, there is no truth-value assignment on which the antecedent of $(\dots (\mathbf{P}_1 \& \mathbf{P}_2) \& \dots \& \mathbf{P}_n) \supset \mathbf{Q}$ has the truth-value \mathbf{T} and the consequent has the truth-value \mathbf{F} . It follows that there is no truth-value assignment on which the corresponding material conditional is false, so it is truth-functionally true.

Assume that $(\dots (\mathbf{P}_1 \& \mathbf{P}_2) \& \dots \& \mathbf{P}_n) \supset \mathbf{Q}$ is truth-functionally true. Then there is no truth-value assignment on which the antecedent is true and the consequent false. But the iterated conjunction is true if and only if the sentences $\mathbf{P}_1, \dots, \mathbf{P}_n$ are all true. So there is no truth-value assignment on which $\mathbf{P}_1, \dots, \mathbf{P}_n$ are all true and \mathbf{Q} is false; hence the argument is truth-functionally valid.

c. No. For example, $\{A \supset B\} \vDash \sim A \vee B$. But $\{A \supset B\}$ does not entail $\sim A$, nor does it entail B .

Section 3.6E

1.a. If $\{\sim \mathbf{P}\}$ is truth-functionally inconsistent, then there is no truth-value assignment on which $\sim \mathbf{P}$ is true (since $\sim \mathbf{P}$ is the only member of its unit set). But then $\sim \mathbf{P}$ is false on every truth-value assignment, so \mathbf{P} is true on every truth-value assignment and is truth-functionally true.

c. If $\Gamma \cup \{\sim \mathbf{P}\}$ is truth-functionally inconsistent, then there is no truth-value assignment on which every member of $\Gamma \cup \{\sim \mathbf{P}\}$ is true. But $\sim \mathbf{P}$ is true on a truth-value assignment if and only if \mathbf{P} is false on that assignment. Hence

there is no truth-value assignment on which every member of Γ is true and \mathbf{P} is false. Hence $\Gamma \vDash \mathbf{P}$.

2.a. \mathbf{P} is truth-functionally true if and only if the set $\{\sim \mathbf{P}\}$ is truth-functionally inconsistent. But $\{\sim \mathbf{P}\}$ is the same set as $\emptyset \cup \{\sim \mathbf{P}\}$. So \mathbf{P} is truth-functionally true if and only if $\emptyset \cup \{\sim \mathbf{P}\}$ is truth-functionally inconsistent. But we have already seen, by previous results, that $\emptyset \cup \{\sim \mathbf{P}\}$ is truth-functionally inconsistent if and only if $\emptyset \vDash \mathbf{P}$. Hence \mathbf{P} is truth-functionally true if and only if $\emptyset \vDash \mathbf{P}$.

c. Assume that Γ is truth-functionally inconsistent. Then there is no truth-value assignment on which every member of Γ is true. Let \mathbf{P} be an *arbitrarily* selected sentence of SL . Then there is no truth-value assignment on which every member of Γ is true and \mathbf{P} false since there is no truth-value assignment on which every member of Γ is true. Hence $\Gamma \vDash \mathbf{P}$.

3.a. Let Γ be a truth-functionally consistent set. Then there is at least one truth-value assignment on which every member of Γ is true. But \mathbf{P} is also true on such an assignment since a truth-functionally true sentence is true on every truth-value assignment. Hence on at least one truth-value assignment every member of $\Gamma \cup \{\mathbf{P}\}$ is true; so the set is truth-functionally consistent.

4.a. \mathbf{P} is either true or false on each truth-value assignment. On any assignment on which \mathbf{P} is true, \mathbf{Q} is true (because $\{\mathbf{P}\} \vDash \mathbf{Q}$) and so $\mathbf{Q} \vee \mathbf{R}$ is true. On any assignment on which \mathbf{P} is false, $\sim \mathbf{P}$ is true, \mathbf{R} is therefore also true (because $\{\sim \mathbf{P}\} \vDash \mathbf{R}$), and so $\mathbf{Q} \vee \mathbf{R}$ is true as well. Either way, then, $\mathbf{Q} \vee \mathbf{R}$ is true—so the sentence is truth-functionally true.

c. Assume that every member of $\Gamma \cup \Gamma'$ is true on some truth-value assignment. Then every member of Γ is true, and so \mathbf{P} is true (because $\Gamma \vDash \mathbf{P}$). Every member of Γ' is also true, and so \mathbf{Q} is true (because $\Gamma' \vDash \mathbf{Q}$). Therefore $\mathbf{P} \& \mathbf{Q}$ is true. So $\Gamma \cup \Gamma' \vDash \mathbf{P} \& \mathbf{Q}$.